

# Communicating Transactions

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## **Transactions in Distributed Systems**

Software Transactional Memory (STM) is a relatively recent mechanism that makes concurrent programming *tractable*, *more modular and composable* [4].

Software Transactions, just like database transactions, are:

- Atomic
- Consistent
- Isolated
- Durable

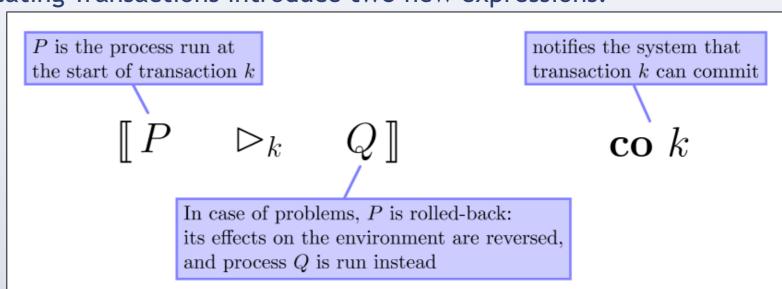


Communication is crucial in distributed systems, such as web service mash-ups. The isolation requirement prevents communication - we can't use traditional transactions!

Communicating transactions drop the I but retain the A, C and D. This makes distributed check-pointing and automated error recovery easy.

#### **The Construct**

Communicating Transactions introduce **two** new expressions:



The system manages communicating transactions through *signals*:

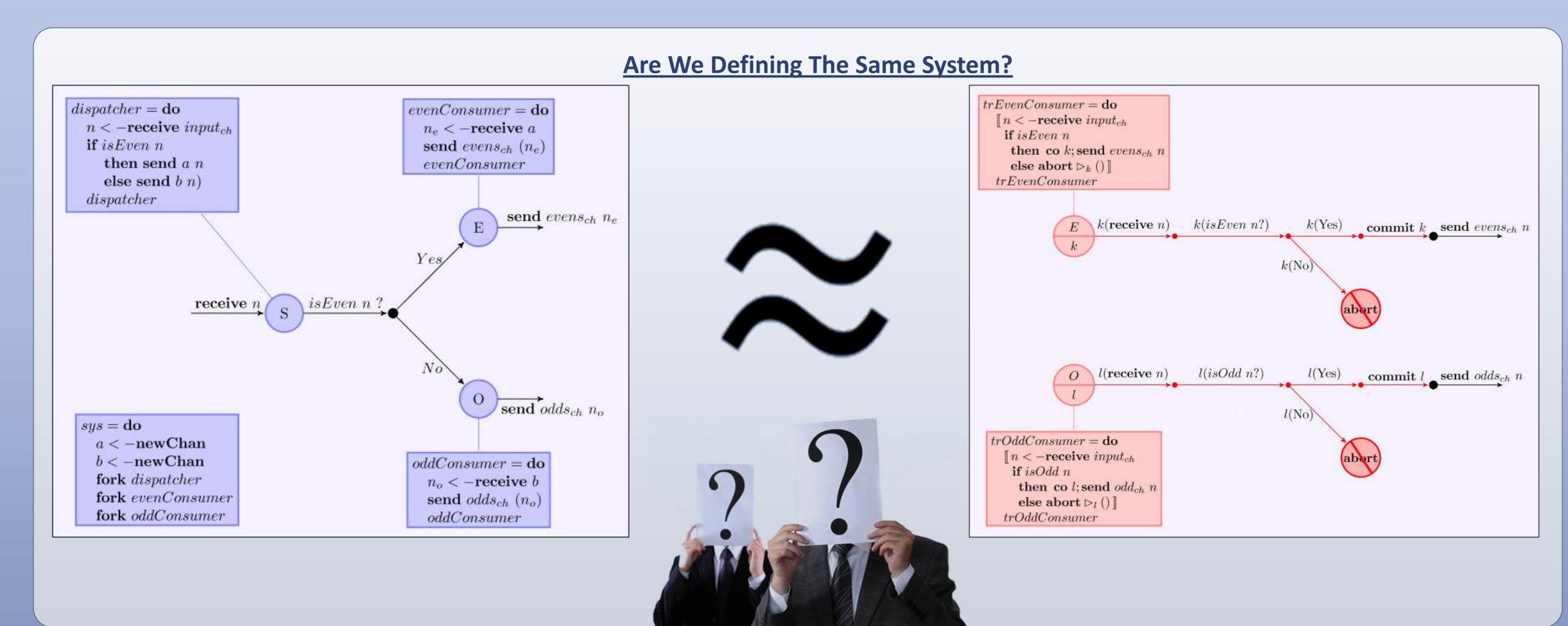
→ Commit: finalizes the transaction.

→ Abort: undoes all actions performed by P and their

effect on the environment since the transaction started.

lets a process join the transaction. Communication is

restricted to processes in the same transaction!



→ Embed:

#### **Leader Election and Nesting** $C_i$ is a configuration where i participants decide to be leaders : undecided : follower : leader $k_1k_2k_3(\tau)$ $k_1k_2k_3(\tau)$ co $k_1$ , co $k_2$ , co $k_3$ $lead(x) = \mathbf{do}$ $a \leftarrow \mathbf{receive} \ c$ $follow(x) = \mathbf{do}$ $b \leftarrow \mathbf{receive} \ c$ $a \leftarrow \mathbf{newChan}$ $y \leftarrow \mathbf{receive} \ a$ $client(v_1) \parallel client(v_2) \parallel client(v_3)$ send c a $z \leftarrow \mathbf{receive}\ b$ send a(x,z)send a xclient(x) =send b(x,y) $(y,z) \leftarrow \mathbf{receive} \ a$ $\llbracket lead(x) \oplus follow(x) \rhd_k 3wR(x) \rrbracket$ co k; (y, z) $\mathbf{co}\ k;\ (y,z)$ Three-way Rendezvous: each participant wants to send his own value and receive the value of the other two participants atomically

### **Research Questions**

- What equivalences can we devise for systems using communicating transactions?
- Can Communicating Transactions be implemented efficiently?
- Are there any useful programming patterns?
- Can we leverage Haskell's type system to create a CT monad?
- How would CT interact with Software Transactional Memory?

#### **Current Status**

- Developed TCML, a transactional version of Concurrent ML, which provides a simple setting to study communicating transactions
- Implementing an interpreter for TCML
- Working on efficient transaction scheduling
- Investigating declarative, Prolog-like patterns for concurrent programming

#### **References and Contact**

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- [4] T. Harris, S. Marlow, S. P. Jones, M. Herlihy: Composable Memory Transactions, PPoPP'05
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