

# Abstraction-Driven Network Verification and Design (a personal odyssey)

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# It started in 2004

- A sabbatical at CMU
  - Joined a collaborative project with AT&T Labs
  - Goal: To reverse engineer the routing designs of 100s of production networks and find ways to detect errors early and minimize outages due to routing loops and blackholes
  - We were given only router configuration files

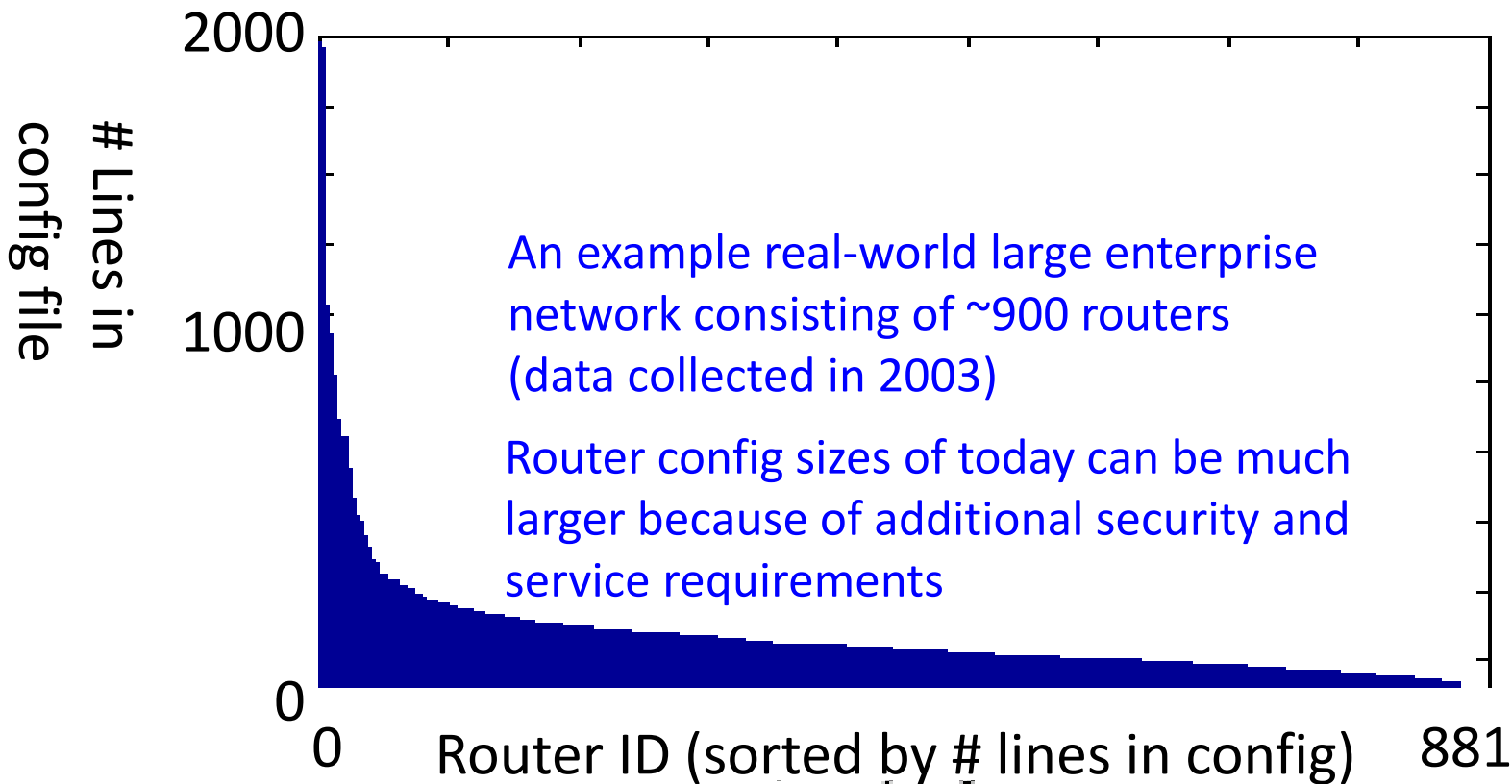
# Excerpts of a Router Config

```
interface Ethernet0
  ip address 6.2.5.14 255.255.255.128
interface Serial1/0.5 point-to-point
  ip address 6.2.2.85 255.255.255.252
  ip access-group 143 in
  frame-relay interface-dlci 28

router ospf 64
  redistribute connected subnets
  redistribute bgp 64780 metric 1 subnets
  network 66.251.75.128 0.0.0.127 area 0
router bgp 64780
  redistribute ospf 64 match route-map 8aTzlvBrbaW
  neighbor 66.253.160.68 remote-as 12762
  neighbor 66.253.160.68 distribute-list 4 in
...
```

```
...
access-list 143 deny 1.1.0.0/16
access-list 143 permit any
route-map 8aTzlvBrbaW deny 10
  match ip address 4
route-map 8aTzlvBrbaW permit 20
  match ip address 7
ip route 10.2.2.1/16 10.2.1.7
...
```

# Lots of Configuration Files



# A Reverse-Engineering Methodology

[Maltz et al, Sigcomm'04]

Configuration files

Find links

Construct logical IP Topology

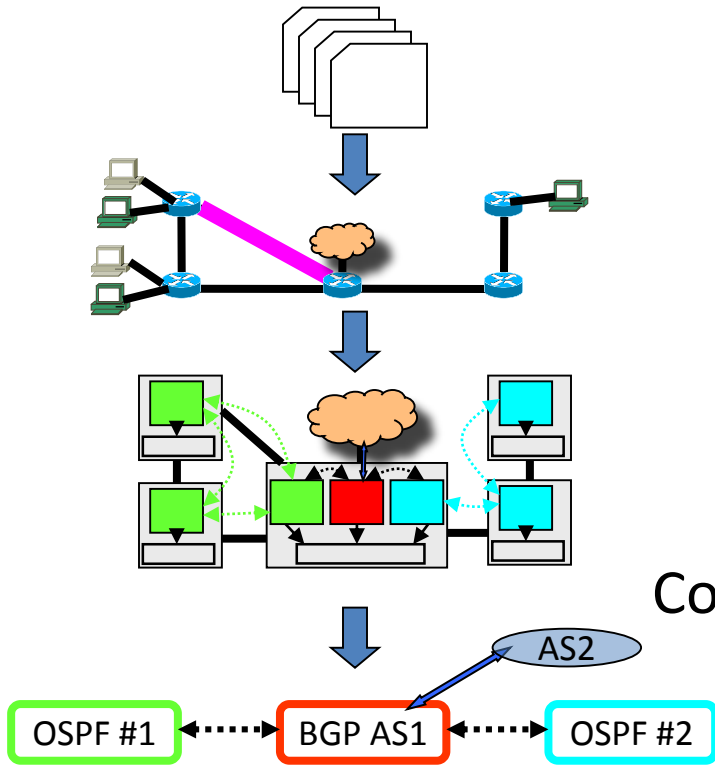
Find adjacent routing processes

Construct Routing Process Graph

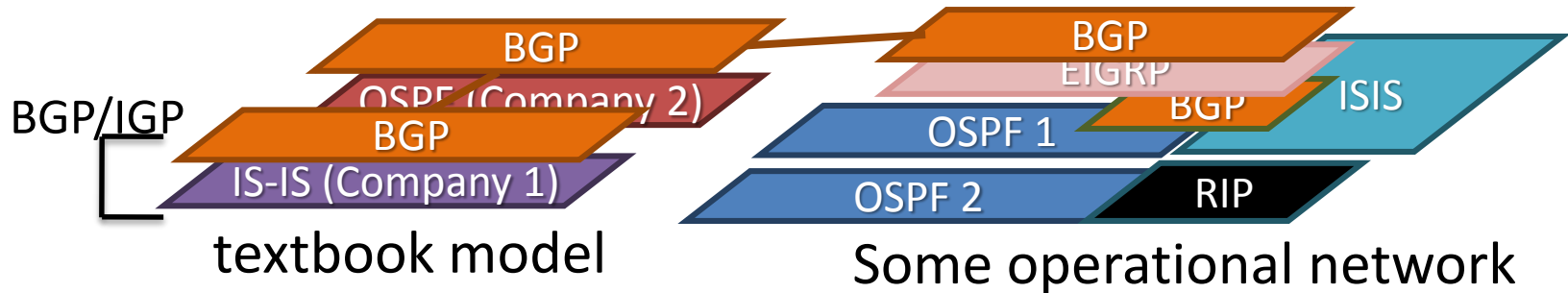
Condense adjacent routing processes into

Routing Instances

This abstraction is key to a scalable solution

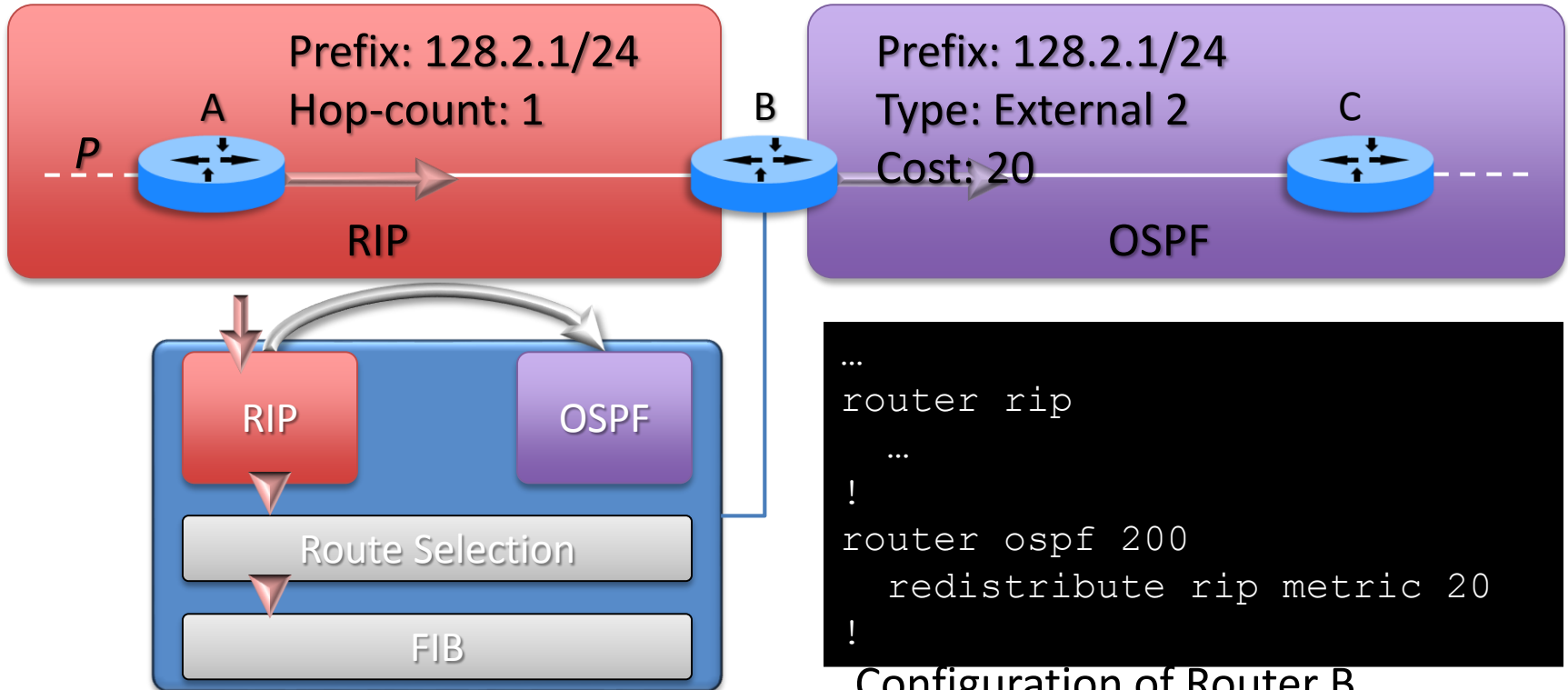


# Diverse Designs beyond Textbook Model



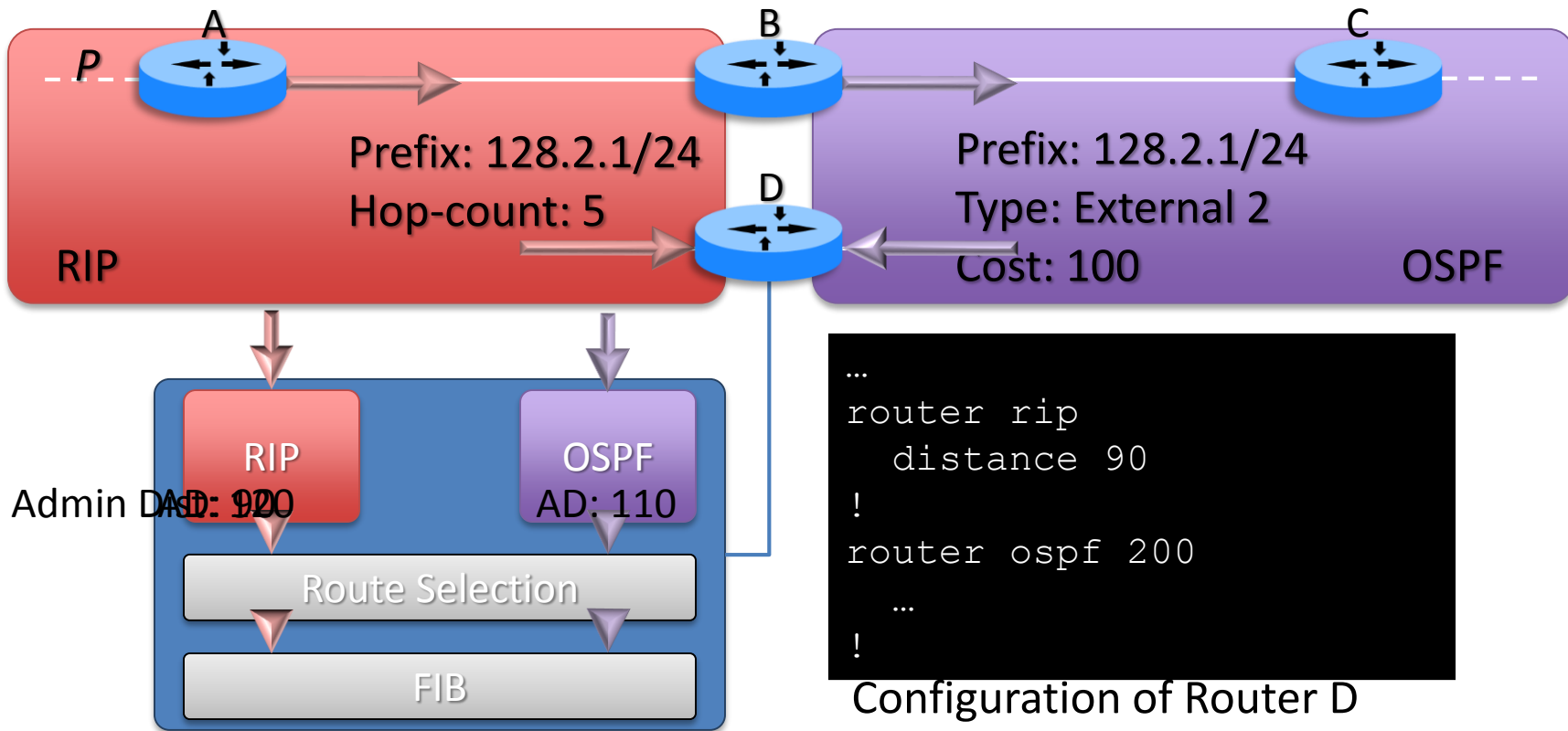
A network can have many routing instances and their interaction will impact routing safety

# Route Redistribution



Configuration of Router B

# Route Selection

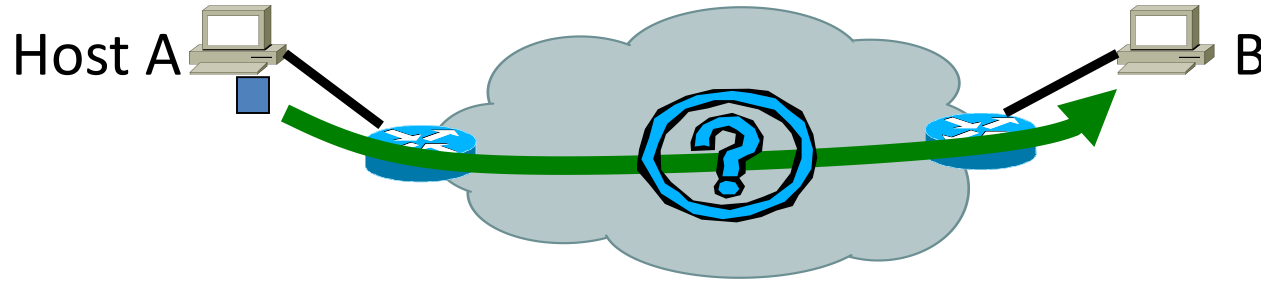




# Impact on Routing Theory & Practice

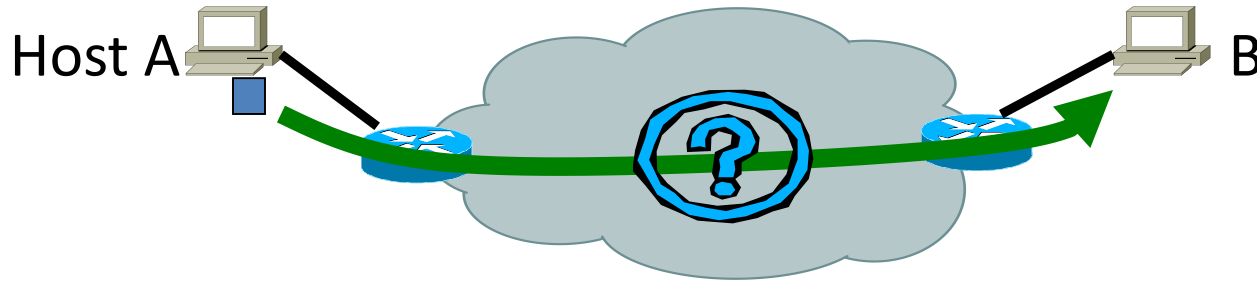
- **Understanding current design and ensuring safety**  
[Le et al, ICNP'07, Sigcomm'08, CoNext'10] [Benson et al., IMC'09, NSDI'09] [Alim and Griffin, CoNext'11] [Sun et al, CoNext'12], etc.
- **Clean slate design of route redistribution**  
[Le et al, Sigcomm'10]
- **Routing Reconfiguration**  
[Vanbever et al, Sigcomm'12] [Vissicchio et al, Infocom'14] etc.
- **Co-existence of multiple control planes (including SDN)**  
[Volpano et al, HotSDN'14] [Vissicchio et al, Infocom'15] etc.

# Next goal: Predicting Reachability



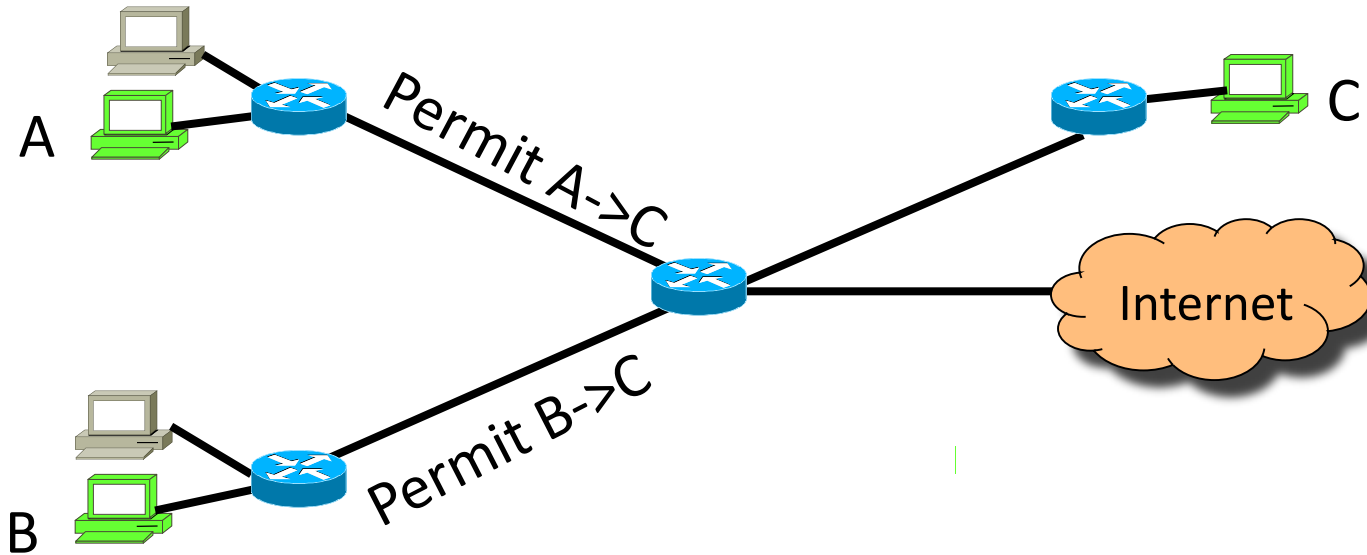
- Reachability depends **not just** on topology
  - Routing protocols, packet filters, and middleboxes
- Predicting reachability is key to network **security** and **resilience**

# State of the Art at the time



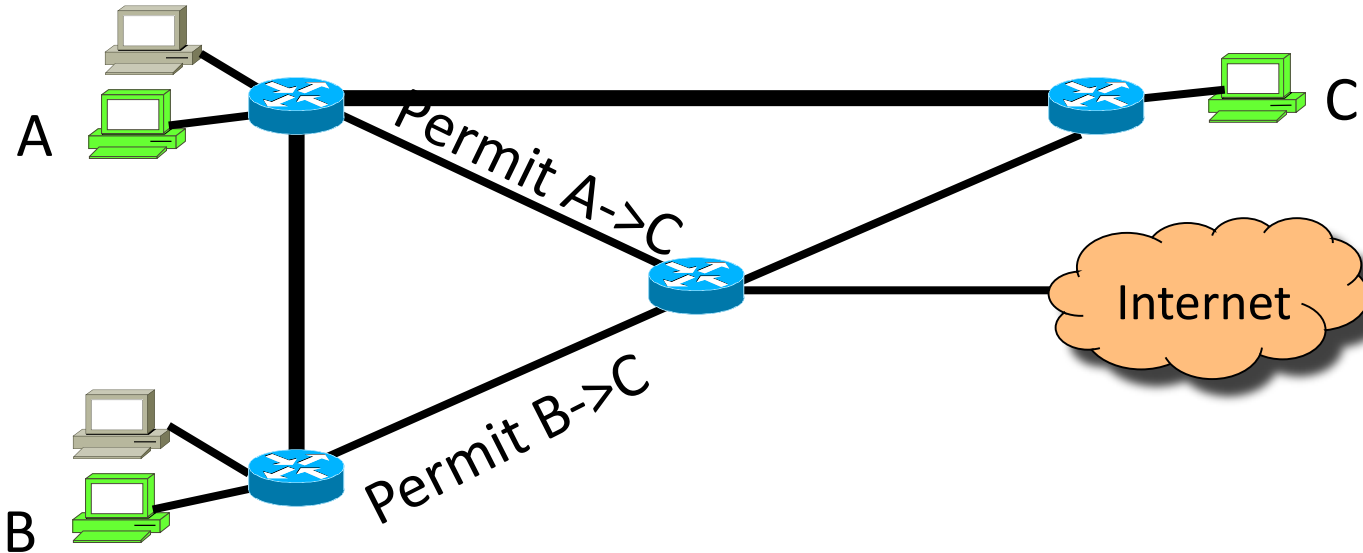
- Build the network and try it
- Dynamic probing (`ping` and `traceroute`) used to troubleshoot reachability problems

# Reachability Example



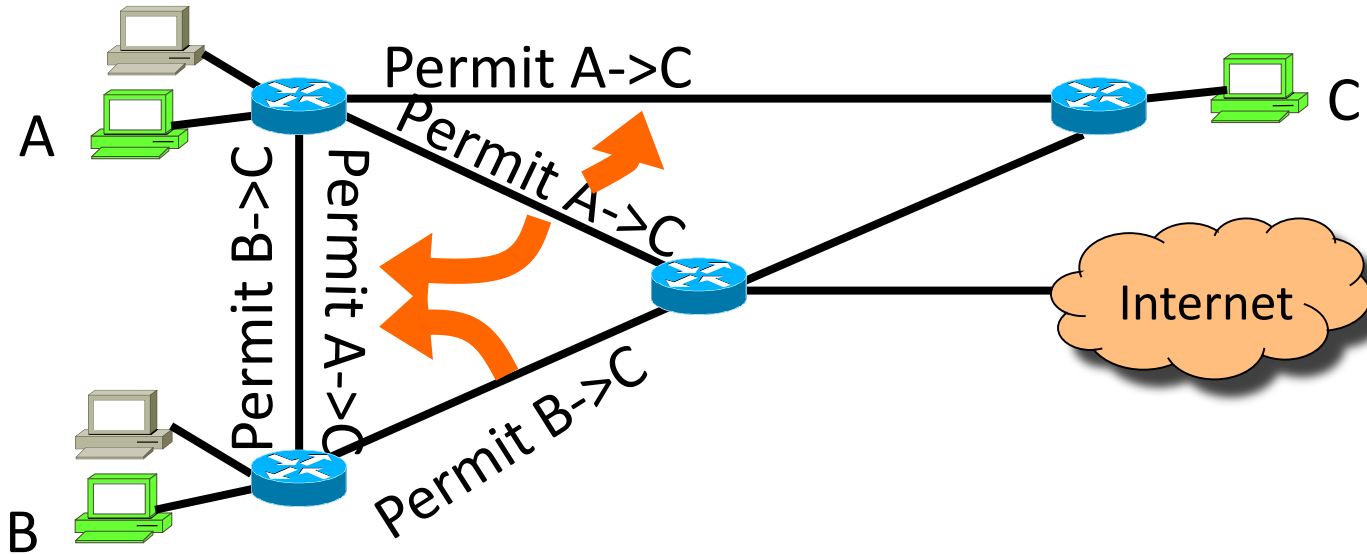
- Enterprise with two remote offices
- Only A&B should be able to talk to server C

# Reachability Example



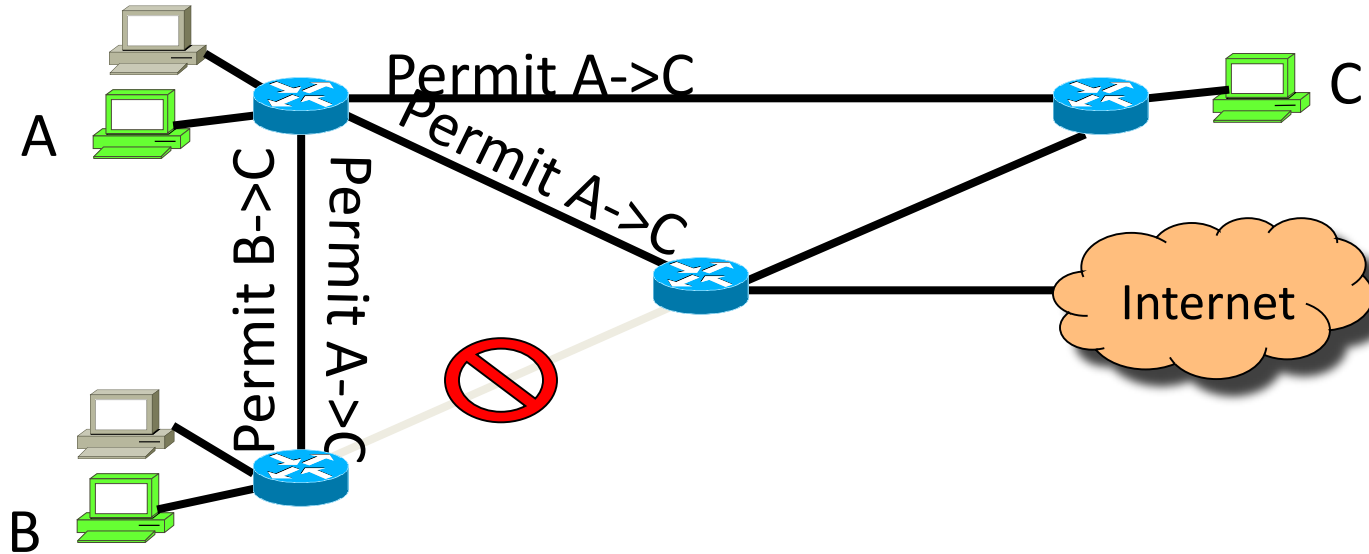
- Network designers add two links for robustness
- Configure routing protocols to use new links in failure

# Reachability Example

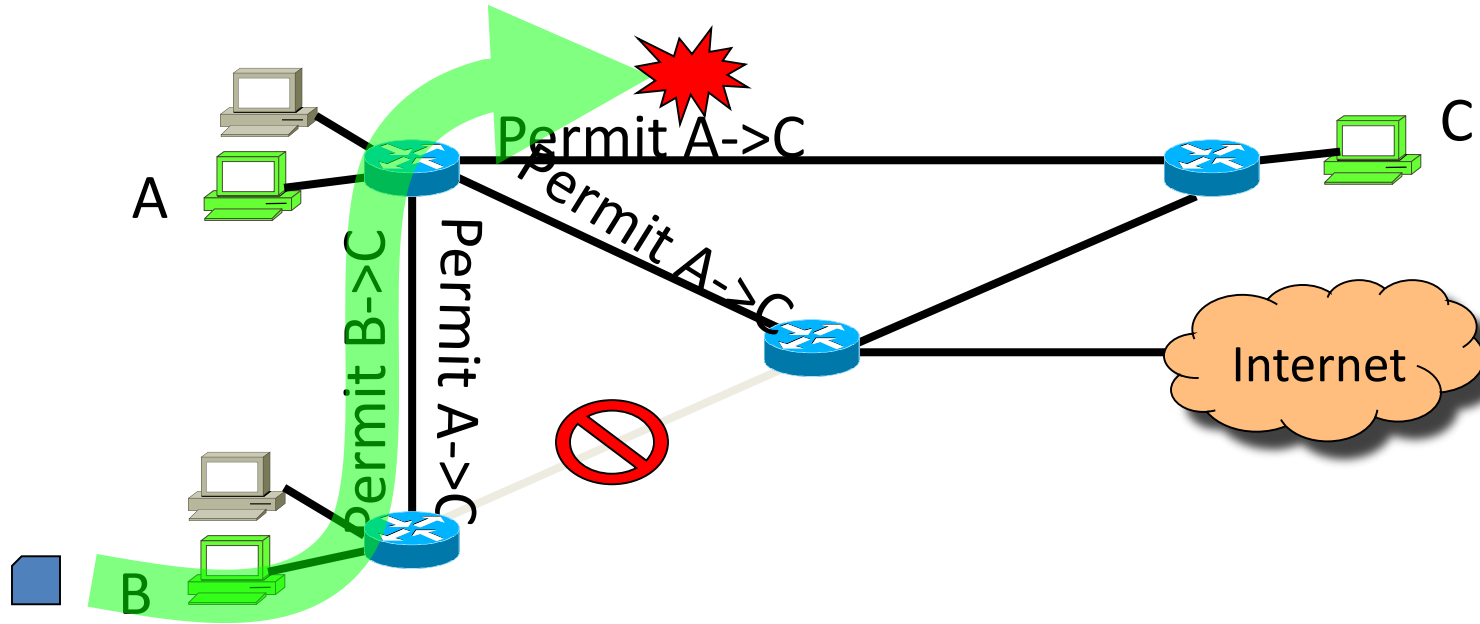


- Designers apply packet filters to new links

# Reachability Example



# Reachability Example



- Packets from B->C dropped!
- Testing under normal conditions won't find this error!



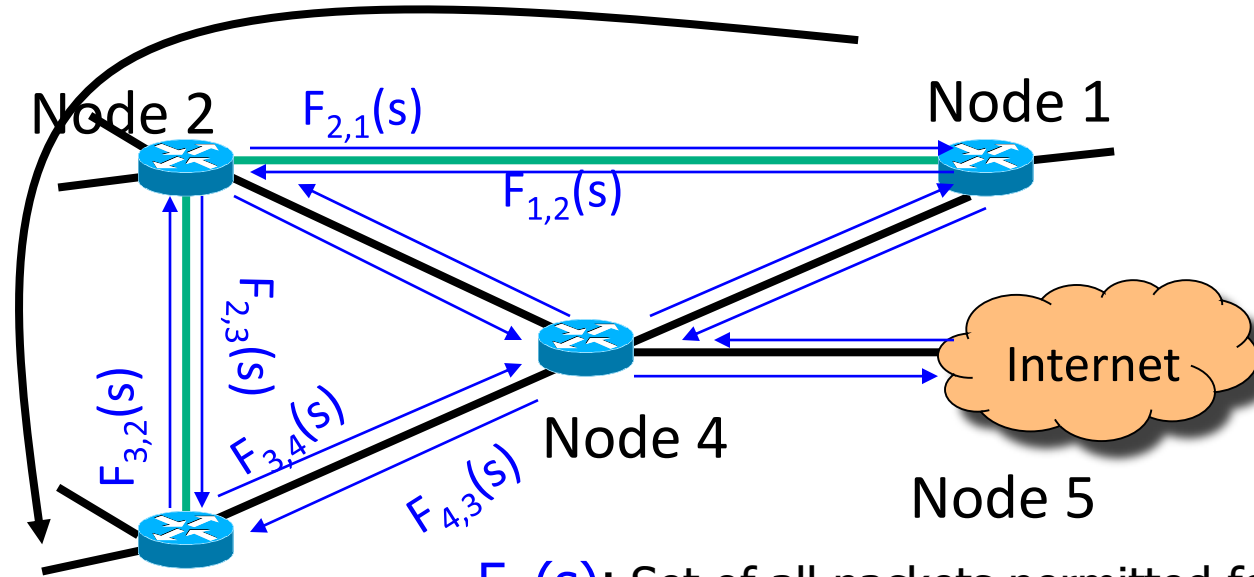
# The Reachability Set abstraction

[Xie et al, Infocom'05]

## Set of all packets permitted from one node to another

- Model **packet filters** naturally  
e.g., "Permit A->C" rule defined on link from node  $u$  to  $v$ :  
 $F_{u,v} = \{\text{packet } p \mid p.\text{src\_addr} = A, p.\text{dst\_addr} = C\}$
- Effect of **routing protocols** added as dynamic destination address based packet filters
  - when network is in forwarding state  $s$ ,  
 $F_{u,v}(s) = F_{u,v} \cap \{\text{all packets } u \text{ would forward to } v \text{ at } s\}$
- **Packet transformation** as generalized inverse function that maps a set of packets to another set of packets

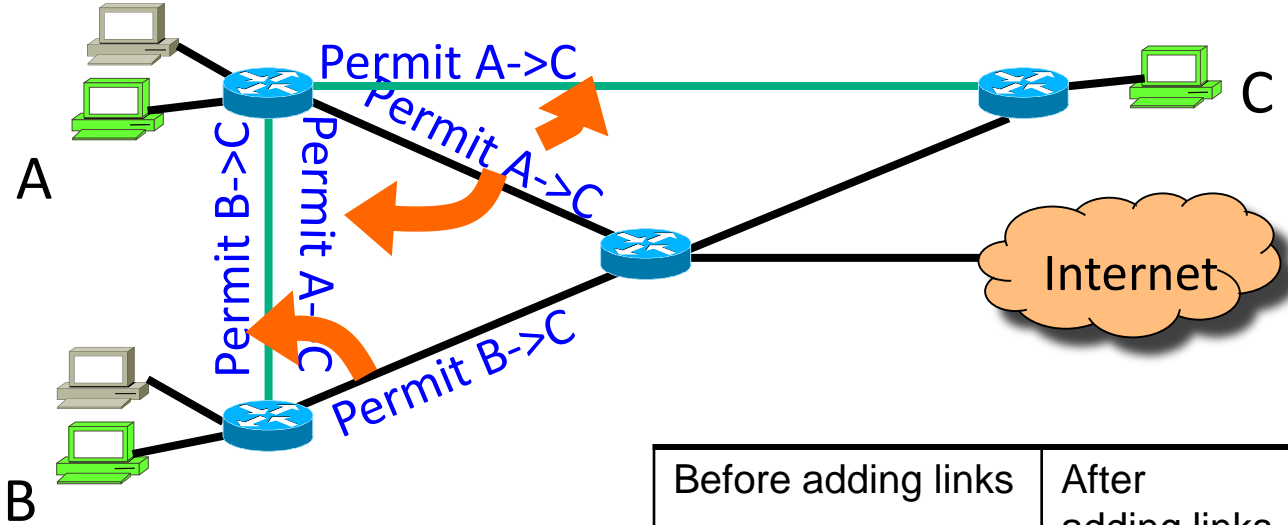
# Reachability Analysis Graph



$F_{i,j}(s)$ : Set of all packets permitted for link  $i \rightarrow j$  at network state  $s$

Reachable Set over directed path  $1 \rightarrow 2 \rightarrow 3 = F_{1,2}(s) \cap F_{2,3}(s)$

# Let's revisit that reachability example



	Before adding links	After adding links
$RS_{B \rightarrow C}^L$	Permit B->C	$(\text{Permit B->C}) \cap (\text{Permit A->C}) = \emptyset !$

# Recent Advances in Static Network Analysis

- Boolean satisfiability formulation  
[Mai et al, Sigcomm'11]
- Header space analysis [Kazemian et al, Sigcomm'12]
- Fast algorithms [Yang and Lam, ICNP'13]

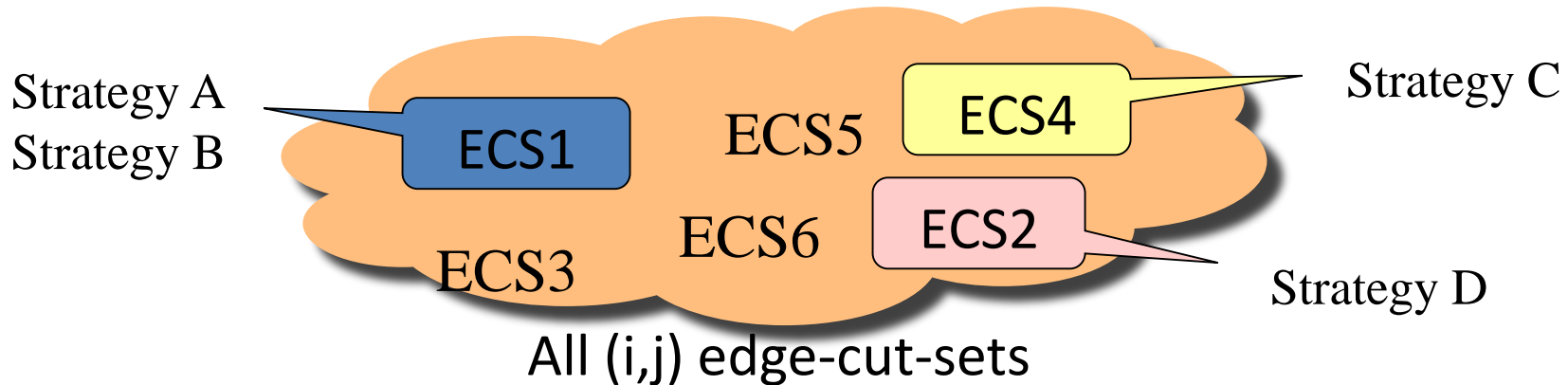
# Next goal: Design automation

[Sung et al CoNext'10]

1. **Abstract** network-wide requirements of a design task
  - Correctness criteria for reachability control modeled as a Reachability Matrix: each cell  $\mathbf{RS}(i,j)$  defines precisely the required reachability set from (virtual) subnet  $i$  to  $j$
2. **Formulate** optimization problems
  - Incorporate resource feasibility constraints (e.g., router capacity for processing packet filter rules)
  - Model explicitly operator strategies (e.g., to deploy a minimum number of filter rules)
3. **Solve** formulated problems
  - Obtain new packet filter placement algorithm

# Automated Packet Filter Placement

- **Intuition:**
  - To achieve  $RS(i, j)$ , same filters must be replicated in an edge-cut-set (ECS) between gateways of the subnets
    - Correctness guaranteed
  - Variety of heuristics possible based on design strategy which chooses particular ECS (minimizing total # of filters, balancing processing load, etc.)



# Some Related Design Efforts

- Integrated design methodology  
[Sun and Xie, CoNext'13]
- Optimizing the “one big switch” abstraction  
[Kang et al, CoNext'13]
- Placement of middleboxes and NFVs  
[Anwer et al, SOSR'15]

# Conclusion

- A huge semantic gap exists between network service objectives and actions of individual protocols and nodes.
  - Software defined networking (SDN) doesn't reduce service objectives, while introducing a new type of nodes and diverse control apps
- Developing higher level abstractions may be key to containing this “curse of many knobs”.
  - E.g., Separation of correctness and performance concerns as in traditional computer programming?