

U-Prove Designated-Verifier Accumulator Revocation Extension

Draft Revision 1

Microsoft Research

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Summary

This document extends the U-Prove Cryptographic Specification [\[UPCS\]](#) by specifying an efficient revocation mechanism based on a dynamic accumulator. This scheme requires a designated verifier that shares the Revocation Authority's private key. Unlike many accumulator schemes based on bilinear pairings, this scheme is built using a prime-order group (like the ones defined in [\[UPCS\]](#)) and is therefore suitable for system that require standard constructions.

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Change history

Version	Date	Description
Draft Revision 1	09/11/2013	Initial draft
	02/26/2014	Added clarification to Figure 3

1 Introduction

This document extends the U-Prove Cryptographic Specification [\[UPCS\]](#) by specifying an efficient revocation mechanism based on a dynamic accumulator. This scheme requires a designated verifier that shares the Revocation Authority's private key. Unlike many accumulator schemes based on bilinear pairings, this scheme is built using a prime-order group (like the ones defined in [\[UPCS\]](#)) and is therefore suitable for system that require standard constructions.

1.1 Notation

In addition to the notation defined in [\[UPCS\]](#), the following notation is used throughout the document.

$a \notin A$ Indicates that element a is not in set A .

The key words "MUST", "MUST NOT", "SHOULD", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [\[RFC 2119\]](#).

1.2 Feature overview

Revocation is an important feature of a credential system; this document specifies a scheme based on a dynamic accumulator using a standard prime order group instead of bilinear pairings.¹

A cryptographic accumulator allows the aggregation of a large set of elements into one constant-size value, and the ability to prove that an element has been aggregated or not in the accumulator.

The accumulator scheme is built on the prime order groups defined in [\[UPCS\]](#); it is therefore possible to use the scheme as a revocation mechanism for U-Prove tokens. The *Revocation Authority* is a new party that manages the revocation list and validates the users' non-revocation proofs. Each token encodes a unique user identifier UID; a non-revocation proof consists of proving that the value UID has not been accumulated in the accumulator. In order to create an efficient non-revocation proof, users periodically obtain revocation witnesses (computed on-demand by the Revocation Authority, or by users as the revocation list is updated); users can then compute in constant time the non-revocation proof.

We detail the DA revocation extension in five steps.

- 1 **Revocation Authority setup:** The Revocation Authority generates its public parameters and secret key, and makes the public parameters available to users.
- 2 **Token issuance:** The user obtains U-Prove tokens encoding her unique identifier UID from the Issuer.
- 3 **Revocation list management:** Periodically, the Revocation Authority updates the revocation accumulator, and the user obtains non-revocation witnesses from the Revocation Authority, or computes them using the revocation list update.
- 4 **Token presentation:** The user presents a U-Prove token to the Verifier, including a non-revocation proof (using the non-revocation witnesses). The Verifier validates the presentation proof.
- 5 **Revocation verification:** The Verifier sends the non-revocation proof to the Revocation Authority that verifies that the undisclosed UID does not appear on the current revocation list.

¹ Many accumulator-based schemes rely on bilinear pairings, common in the cryptographic literature, but not yet popular in industry systems and standards.

2 Protocol specification

2.1 Revocation Authority setup

The Revocation Authority generates its key and parameters as specified in Figure 1. The group G_q MUST match the one specified in the parameters of Issuers that will issue tokens revocable by this Revocation Authority.

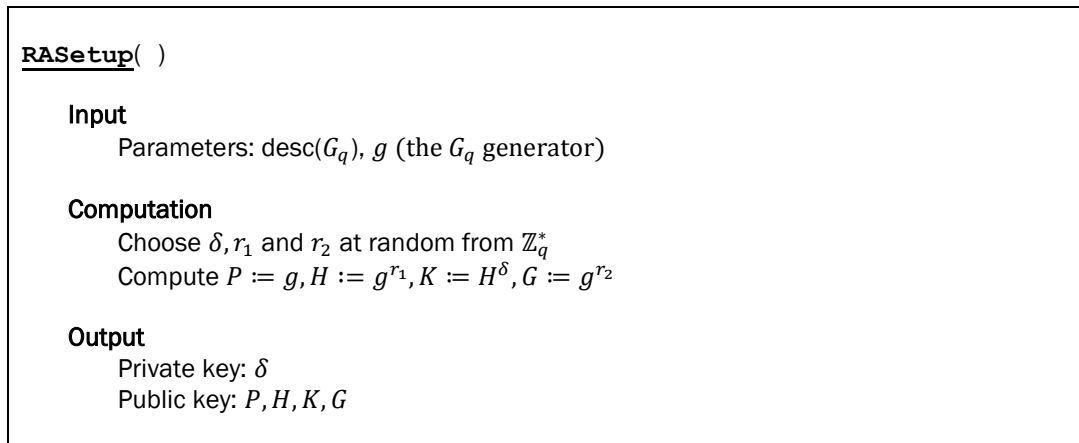


Figure 1: Function RAS_{Setup}

2.2 Token issuance

Token issuance follows the same steps as in [\[UPCS\]](#). One of the attributes, called the *revocation attribute* and denoted x_{id} (where id is an index value between 1 and the number of attributes in the tokens), is reserved for revocation.

2.3 Revocation list management

The Revocation Authority computes the accumulator corresponding to a set of revoked attribute values² (see Figure 2); the accumulator is re-computed when values are added or removed from the revocation list. Users periodically obtain revocation witnesses corresponding to their revocation attribute x_{id} allowing them to create non-revocation proofs. If the revocation list is secret, or for better efficiency, the witnesses are computed by the Revocation Authority (see Figure 3); otherwise, the witnesses are computed by users and updated when values are added or removed from the revocation list (see Figure 4).

² Initially, this set can be empty.

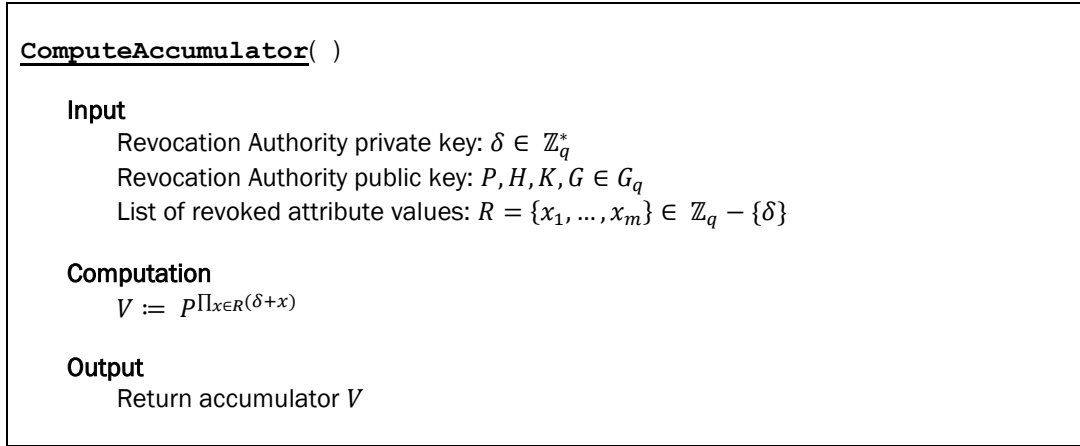


Figure 2: Function ComputeAccumulator

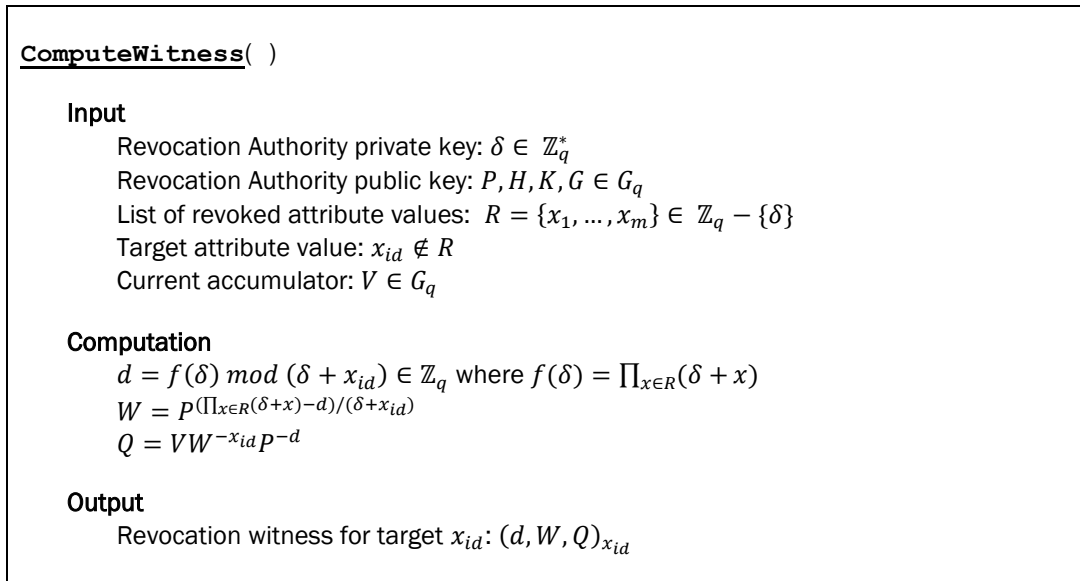


Figure 3: Function ComputeWitness

Note that the computation of d is a polynomial division of polynomial $f(\delta) = \prod_{x \in R} (\delta + x)$ over polynomial $(\delta + x_{id})$ in polynomial ring $\mathbb{Z}_q[\delta]$. As the denominator is a polynomial of degree 1, the result d is a polynomial of degree 0, i.e. a constant, in the polynomial ring. So d can be computed from just the set R and does not depend on δ .

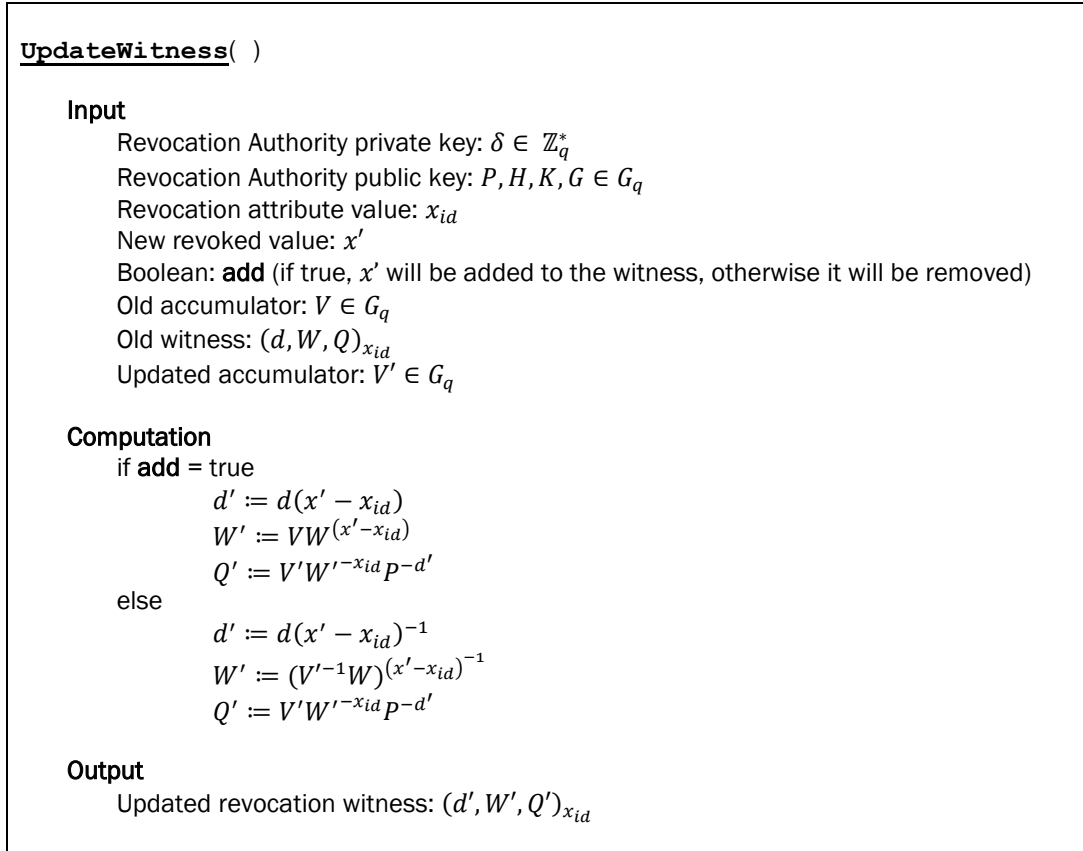


Figure 4: Function UpdateWitness

2.4 Token presentation

The presentation proof is generated according to the needs of the application following [\[UPCS\]](#), but additionally x_{id} is a committed attribute. The (public) output $\tilde{c}_{id} = g^{x_{id}}g_1^{o_{id}}$ and the (private) opening information (x_{id}, o_{id}) , are input to the non-revocation proof generation defined in Figure 5.

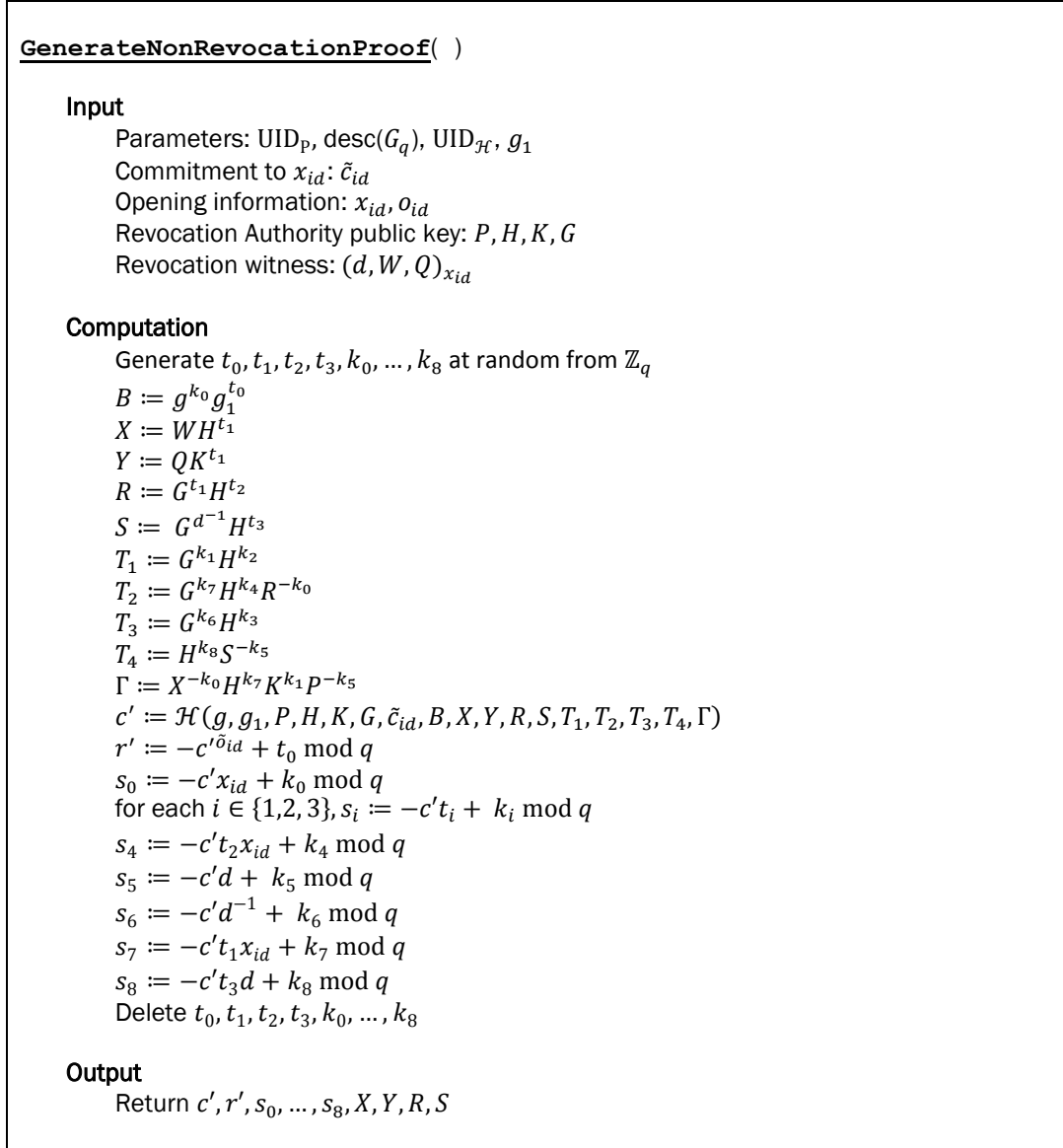


Figure 5: Function GenerateNonRevocationProof.

2.5 Revocation verification

The revocation verification function, defined in Figure 6, is run after the corresponding presentation proof has been successfully verified. Inputs are the commitment \tilde{c}_{id} from the presentation proof and the non-revocation proof. If the presentation and non-revocation proofs are valid, then the verifier has assurance that \tilde{c}_{id} is a valid commitment to the attribute x_{id} and that x_{id} is not in the revocation list.

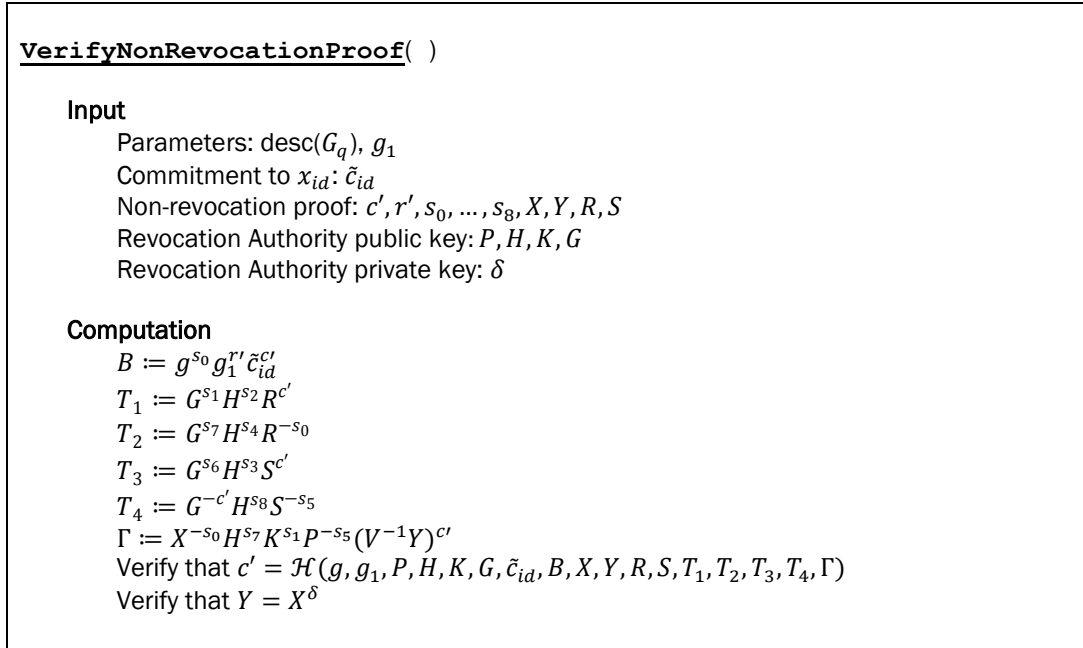


Figure 6: Function VerifyNonRevocationProof

References

- [RFC2119] Scott Bradner. *RFC 2119: Key words for use in RFCs to Indicate Requirement Levels*, 1997. <ftp://ftp.rfc-editor.org/in-notes/rfc2119.txt>.
- [UPCS] Christian Paquin, Greg Zaverucha. *U-Prove Cryptographic Specification V1.1 (Revision 2)*. Microsoft, April 2013. <http://www.microsoft.com/u-prove>.